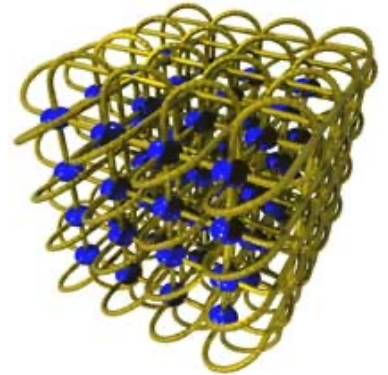


January, 2003



Super-scalable Algorithms

Next Generation Supercomputing on
100,000 and more Processors

Christian Engelmann

Overview

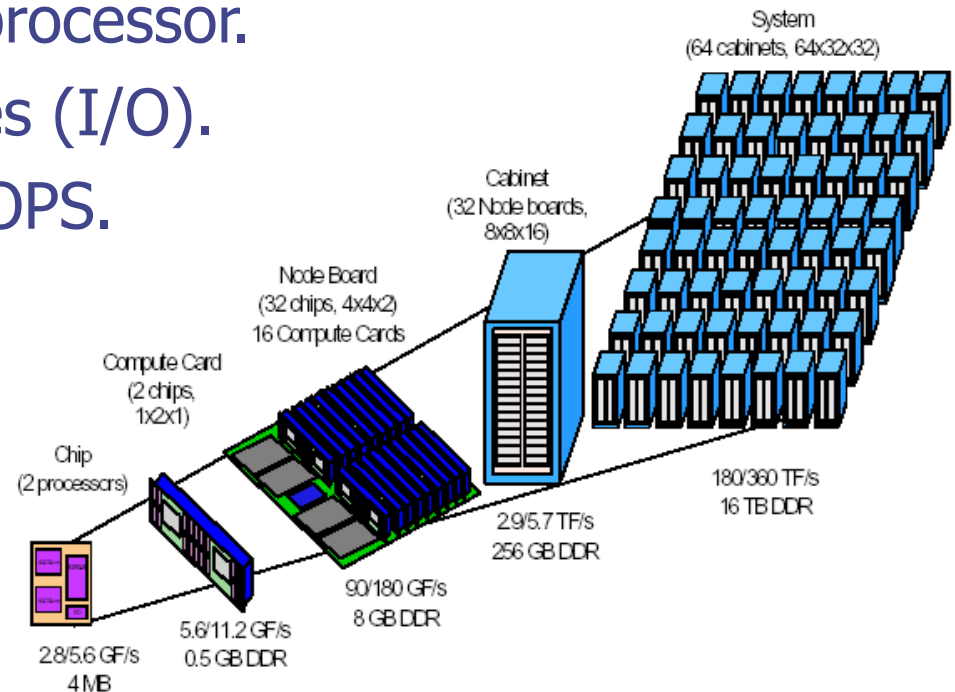
- ◆ Super-scale architectures.
- ◆ Scalability and fault-tolerance issues.
- ◆ Cellular algorithms theory.
- ◆ ORNL/IBM collaboration.
- ◆ IBM BlueGene\L emulators.
- ◆ ORNL cellular architecture simulator.
- ◆ Super-scalable algorithms.
- ◆ Super-scalable diskless checkpointing.
- ◆ Conclusions and ideas for the future.

Super-scale Architectures

- ◆ Current tera-scale supercomputers have up to 10,000 processors.
- ◆ Next generation peta-scale systems will have 100,000 processors and more.
- ◆ Such machines may easily scale up to 1,000,000 processors in the next decade.
- ◆ IBM currently builds the BlueGene\L at Lawrence Livermore National Laboratory.

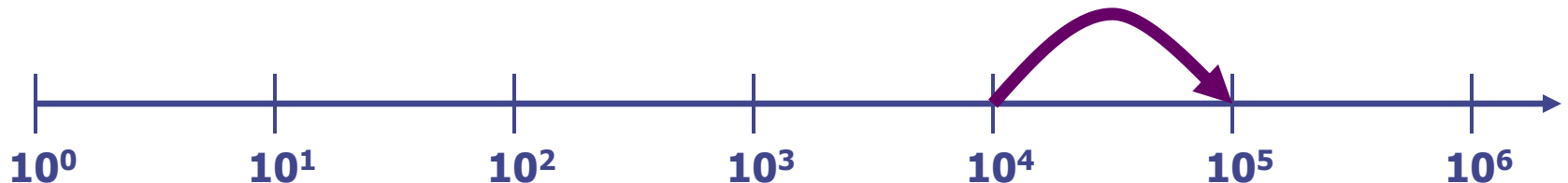
IBM BlueGene/L at LLNL

- ◆ Up to 64K diskless nodes with 2 processors per node.
- ◆ Only 256MB RAM per processor.
- ◆ Additional service nodes (I/O).
- ◆ Estimated 360 Tera FLOPS.
- ◆ Over 150k processors.
- ◆ Global tree network.
- ◆ 3-D torus network.
- ◆ Gigabit Ethernet.
- ◆ Operational in 2005.



Scalability Issues

- ◆ How to make use of 100,000 processors?
- ◆ System scale jumps by a magnitude.
- ◆ Current algorithms do not scale well on existing 10,000-processor systems.
- ◆ Next generation peta-scale systems are useless if efficiency drops by a magnitude.



Fault-tolerance Issues

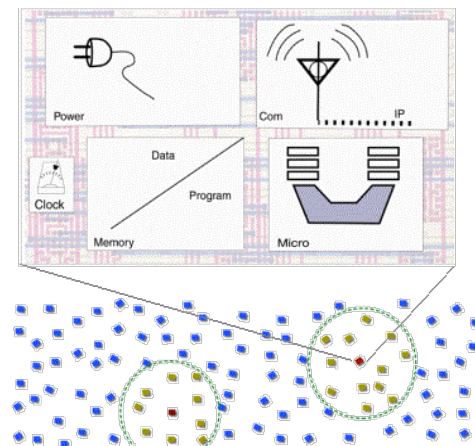
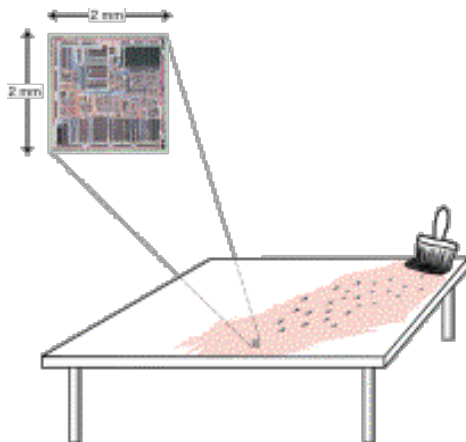
- ◆ How to survive on 100,000 processors?
- ◆ Failure rate grows with the system size.
- ◆ Mean time between failures may be a few hours or just a few minutes.
- ◆ Current solutions for fault-tolerance rely on checkpoint/restart mechanisms.
- ◆ Checkpointing 100,000 processors to central stable storage is not feasible anymore.

Cellular Algorithms Theory

- ◆ Processes have only limited knowledge mostly about other processes in their neighborhood.
- ◆ Application is composed of local algorithms.
- ◆ Less inter-process dependencies, e.g not everyone needs to know when a process dies.
- ◆ Peer-to-peer communication with overlapping neighborhoods promotes scalability.

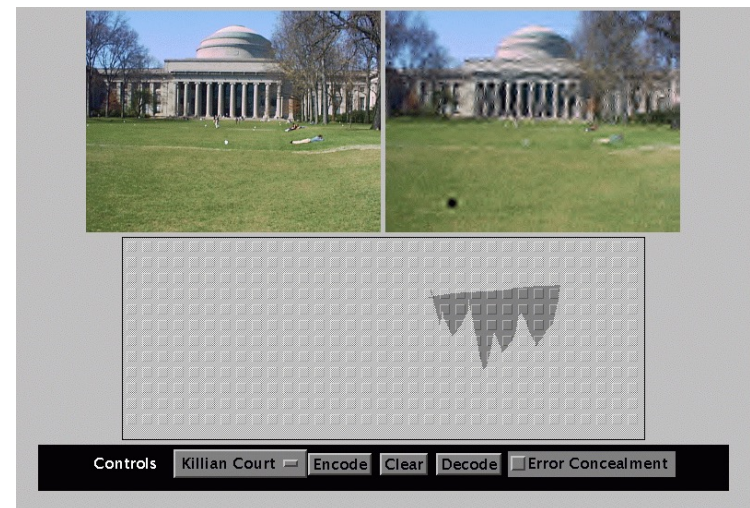
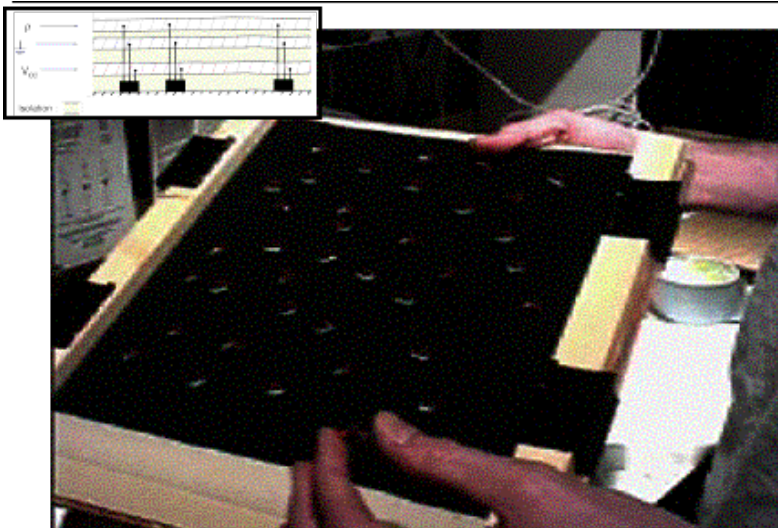
Paintable Computing at MIT

- ◆ In the future embedded computers with a radio device get as small as a pigment.
- ◆ Supercomputers can be easily assembled by painting a wall of embedded computers.
- ◆ Applications are driven by cellular algorithms.



Paintable Computing at MIT

- ◆ 2002 Ph.D. Thesis using a pushpin board.
- ◆ Applications:
 - Distributed audio stream storage.
 - Fault-tolerant holistic data storage.



ORNL/IBM Collaboration

- ◆ Development of biology and material science applications for super-scale systems.
- ◆ Exploration of super-scalable algorithms.
 - Natural fault-tolerance.
 - Scale invariance.
- ◆ Focus on test and demonstration tool.
- ◆ Get scientists to think about scalability and fault-tolerance in super-scale systems!

ORNL Research Group

- ◆ Al Geist (PI).
- ◆ Christian Engelmann (simulator).
- ◆ Kasidid Chanchio (global max problem).
- ◆ Ryan Adamson (async. multigrid).
- ◆ Bill Shelton (LSMS port).
- ◆ Pratul Agarwal (MD port).

BlueGene\L Emulators

◆ IBM Research:

- Processor emulation with OS in a Linux process.

◆ Caltech:

- MPI trace file analysis for performance prediction.

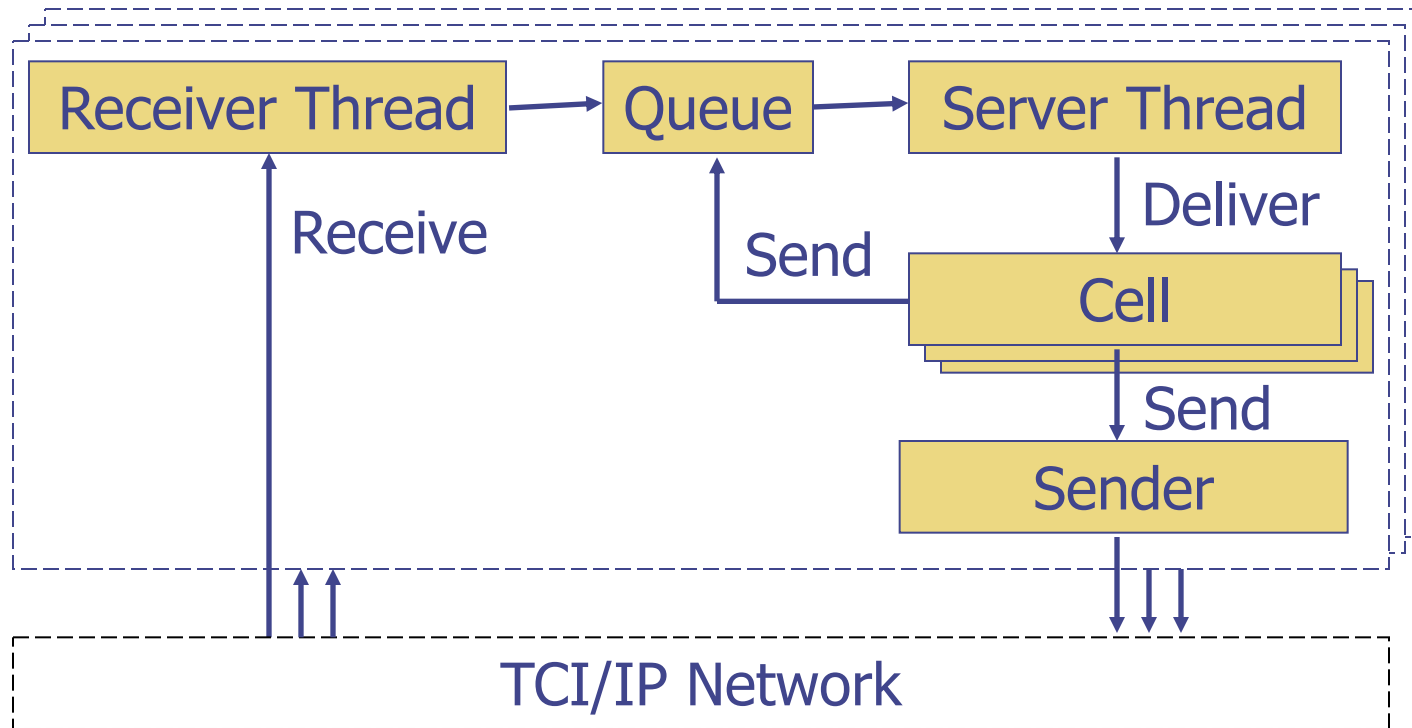
◆ UIUC:

- Object-oriented message driven emulation of logical system architecture in Converse/Charm++.
- Adaptive MPI emulation on top of Charm++.
- Scalability and performance issues in prototypes.
- Emulation fixed on BlueGene\L architecture.

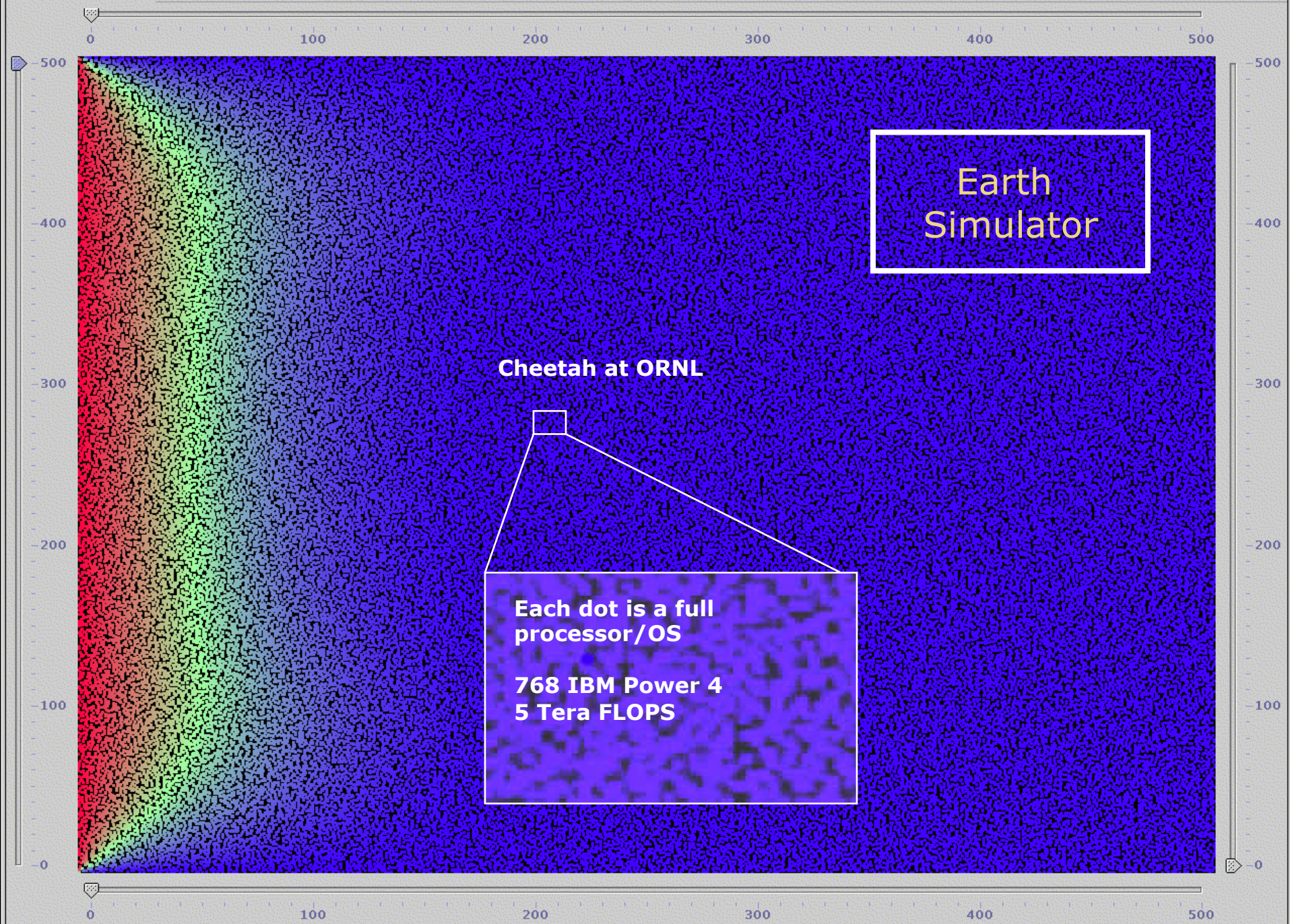
Cellular Architecture Simulator

- ◆ Developed at ORNL in Java with native C and Fortran application support using JNI.
- ◆ Runs as standalone or distributed application.
- ◆ Lightweight framework simulates up to 1,000,000 processes on 9 real processors.
- ◆ Standard and experimental networks:
 - Multi-dimensional mesh/torus.
 - Nearest/Random neighbors.
- ◆ Message driven simulation is not in real-time.
- ◆ Primitive fault-tolerant MPI support.

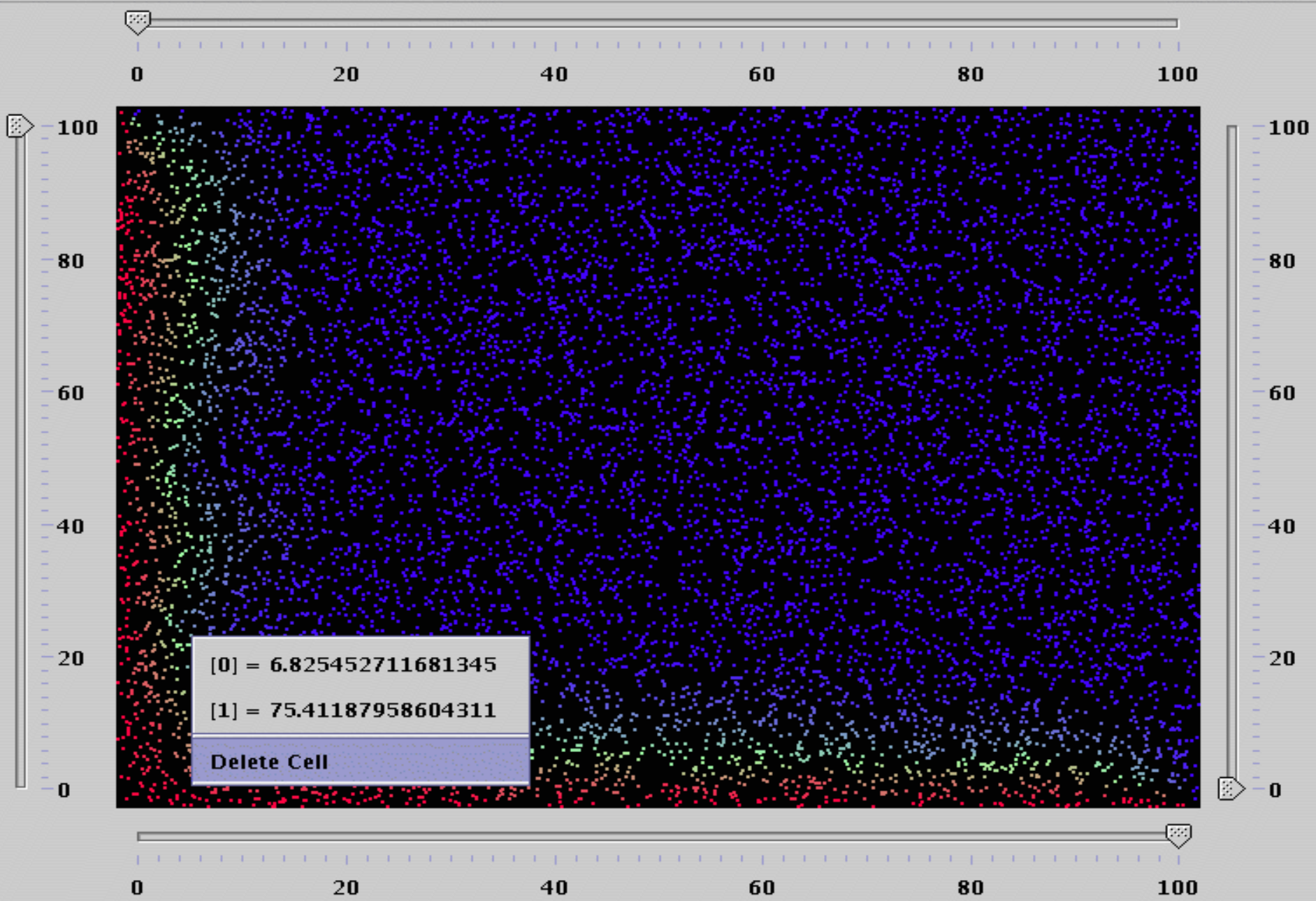
Cellular Architecture Simulator



- ◆ Every cell has its own code, memory and neighbors list.
- ◆ Server hosts cells and initiates the context switch.
- ◆ Cells communicate asynchronously using messages.



System Laplace (Java) Help

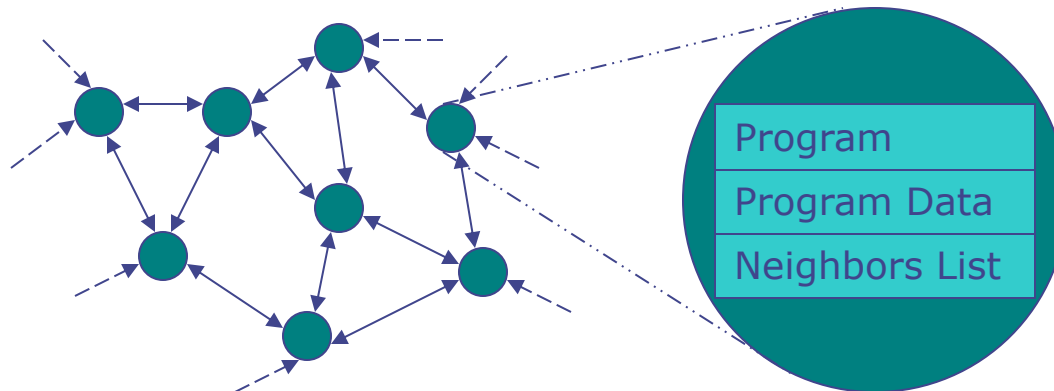


Super-scalable Algorithms

- ◆ Extending the cellular algorithms theory to real world scientific applications.
- ◆ Exploring super-scale properties:
 - Scale invariance
 - Natural fault-tolerance.
- ◆ Gaining experience in programming models for 100,000-processor machines.

Scale invariance

- ◆ Linear scalability.
- ◆ Peer-to-peer communication patterns are based on a small set of neighbor processes.
- ◆ Neighbors are random, far away or nearby.
- ◆ Global application state is composed of many interdependent local neighborhood states.



Natural Fault-tolerance

- ◆ Ability to get the correct answer despite task failures and without checkpointing.
- ◆ May involve redundant computation.
- ◆ 0,1% failure rate (100 of 100,000 processors) is still acceptable with 0,5% redundancy.
- ◆ Failures detected by hardware and ignored or accepted by neighbor processes.
- ◆ Failed processes may be restarted by “inserting” new ones at anytime.

Researched Algorithms

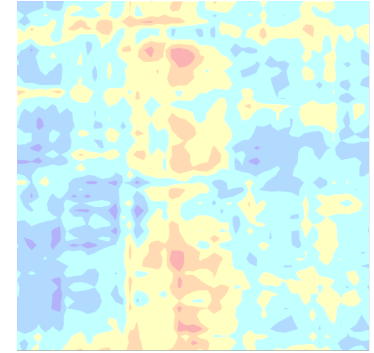
◆ Local information exchange:

- Local peer-to-peer updates of values.
- Mesh-free chaotic relaxation (Laplace/Poisson).
- Finite difference/element methods.
- Dynamic adaptive refinement at runtime.
- Asynchronous multi-grid with controlled or independent updates between different layers.

◆ Global information exchange:

- Global peer-to-peer broadcasts of values.
- Global maximum/optimum search.

Ported Applications



◆ Material Science:

- Magnetism simulation using the locally self-consistent multiple scattering (LSMS) method for understanding the interactions between electrons and atoms in magnetic materials (Bill Shelton).

◆ Computational Biology:

- Molecular dynamics (MD) simulation of biological molecules (DNA sequences) for understanding the protein-DNA interactions (Pratul Agarwal).

Observations

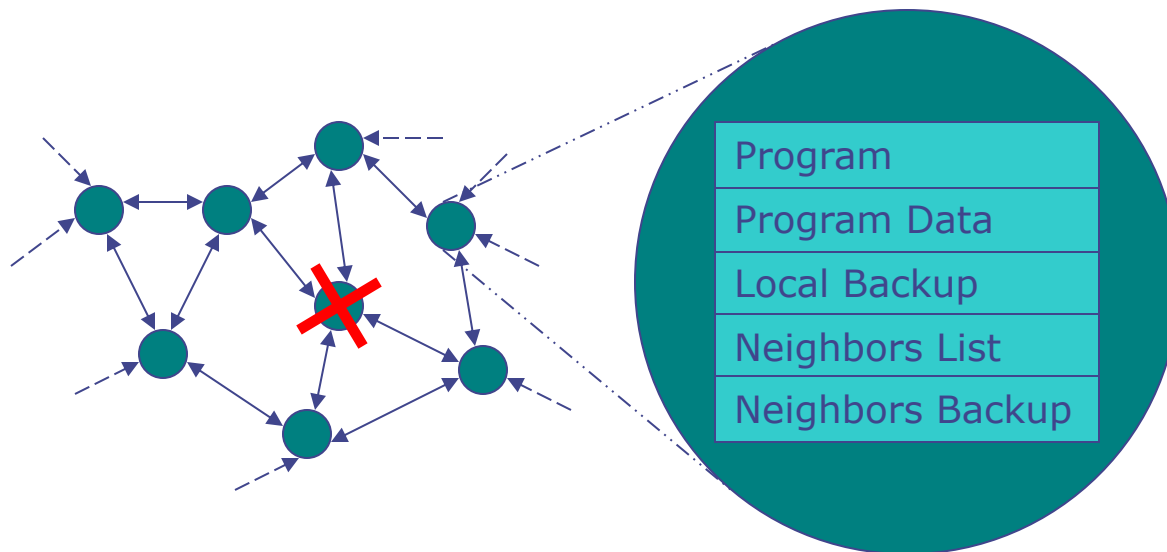
- ◆ Partially non-deterministic algorithm behavior.
- ◆ Unpredictable application running time.
- ◆ Chaotic relaxation does not always converge.
- ◆ No exact replay without full message trace.
- ◆ Communication bound algorithms that require high point-to-point bandwidth.
- ◆ Asynchronous message driven programming model similar to discrete event simulations.
- ◆ Message queues with overwrite.

Super-scalable Fault-tolerance

- ◆ For non-naturally fault tolerant algorithms.
 - ◆ Does it makes sense to restart all 100,000 processors because one failed?
 - ◆ The mean time between failures is likely to be a few hours or just a few minutes.
 - ◆ Traditional centralized checkpointing is limited by bandwidth (bottleneck).
- The failure rate is going to outrun the recovery and the checkpointing rate.

Diskless Checkpointing

- ◆ Decentralized peer-to-peer checkpointing.
- ◆ Processors hold backups of neighbors.
- ◆ Local checkpoint and restart algorithm.
- ◆ Coordination of local checkpoints.



Diskless Checkpointing

- ◆ In case of a failure:
 - Rollback to local memory backup if necessary.
 - Restart from remote memory backup.
- ◆ Encoding semantics, such as RAID, trade off storage size vs. degree of fault tolerance.
- ◆ Very infrequent checkpointing to central stable storage (disk/tape).
- ◆ Checkpoint and application processes may be the same or different.
- ◆ Possible OS support via library/service.

Choosing Neighbors

- ◆ Physically near neighbors:
 - Low latency, fast backup and recovery.
- ◆ Physically far neighbors:
 - Recoverable multiprocessor node failures.
- ◆ Random neighbors:
 - Medium latency and bandwidth.
 - Acceptable backup and recovery time.
- ◆ Optimum: Pseudorandom neighbors based on system communication infrastructure.

Backup Coordination

- ◆ All checkpoints need to be consistent with the global application state.
- ◆ Local states and in-flight messages.
- ◆ No coordination for checkpoints with no communication since last or since start.
- ◆ Coordination techniques:
 - Global synchronization.
 - Local synchronization.

Global Synchronization

- ◆ Global application snapshot (e.g. barrier) at stable global application state.
- ◆ Synchronous backup of all local states.
- ◆ Easy to implement.
- ◆ Synchronizes complete application.
- ◆ Preferred method for communication intensive applications.

Local Synchronization

- ◆ Asynchronous backup of local state and in-flight messages (message logging).
- ◆ Acknowledgements for messages to keep accurate records of in-flight messages.
- ◆ Additional local group communication.
- ◆ Different methods to retrieve missed messages from neighbors.
- ◆ More complicated to implement.
- ◆ Preferred method for less communication intensive applications.

Observations

- ◆ Diskless peer-to-peer checkpointing on super-scale architectures is possible.
- ◆ Synchronization methods have different strengths and weaknesses.
- ◆ Timing, latency and bandwidth data impossible to obtain from simulator.
- ◆ Real-time tests with different applications are needed for further discussion.
- ◆ Final real-world implementation requires super-scalable FT-MPI or PVM.

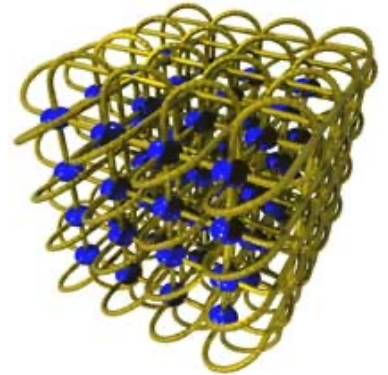
Conclusions

- ◆ Super-scale systems with 100,000 and more processors become reality very soon.
- ◆ Super-scalable algorithms that are scale invariant and naturally fault-tolerant do exist.
- ◆ Diskless peer-to-peer checkpointing provides an alternative to natural fault-tolerance.
- ◆ A lot of research still needs to be done.

Ideas for the Future

- ◆ Research in OS and/or middleware supported super-scale diskless checkpointing.
- ◆ Development of super-scalable fault-tolerant MPI implementation with localized recovery.
- ◆ Development of super-scalable algorithms for specific applications in computational biology, material science, climate research ...

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