## **System-level Virtualization for High Performance Computing**

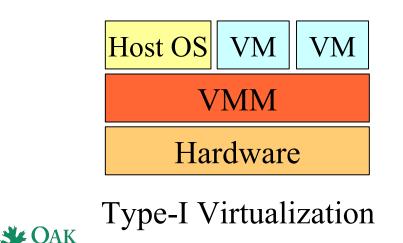
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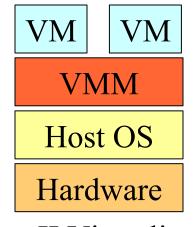


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## **Introduction to Virtualization**

- System-level virtualization studied since the 70's (Goldberg, Popek)
- Key concepts:
  - Virtual Machine (VM), guest OS: complete operating system running in a virtual environment
  - Host OS: operating system running on top of the hardware, interface between the user and the VMM and VMs
  - Virtual Machine Monitor (VMM), Hypervisor: manage VMs (scheduling, hardware access)





Type-II Virtualization

## **The HPC Context**

- The system should not interfere with running application(s)
  - minimize the OS footprint/noise: Catamount, CNK
  - priority to the application for resources access
- Support large-scale systems
  - thousands of distributed components
  - fault tolerance
  - system partitioning: compute nodes vs. service nodes
- Support HPC applications, e.g., MPI applications

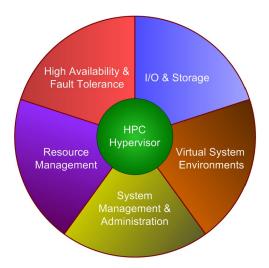


# **Hypervisor for HPC**

- What does it mean?
  - VMM for HPC
    - minimize the VMM memory footprint
    - VMM/HostOS system footprint
      - only the HostOS can execute privileged instructions on the behalf of VMs
      - possible domain context switches
  - Provide a suitable *execution environment* to HPC applications
  - Fault tolerance
  - System Management of VMs/VMMs/HostOSes in large-scale distributed systems
  - I/O & Storage
    - efficient access to resources
    - resource sharing between VMs running on different nodes

### Current virtualization solutions are not suitable for HPC





### **Hypervisor for High-Performance Computing**



## **VMM for High-Performance Computing**

- Minimize the system footprint
  - reduce the default VMM/HostOS memory usage
  - use hardware optimizations (AMD nested pages, hardware IOMMU)
- Minimize the context switches between domains
  - pin Vms/HostOSes/VMM to a core/processor
- Avoid the usage of HostOSes for a direct access to the bare hardware
  - isolation vs. resource sharing
  - "VMM-bypass"

#### Guarantee experiment reproductability

Have the same behavior between different application runs



### **Virtual System Environment**



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## **Virtual System Environment (VSE)**

- Objective: "Adapt the operating system to the application, instead of adapting the application to the OS"
  - The science resides in the applications, not in the operating systems or run-times
  - Goal: application developers should not "port" their application every time they want to use a new execution platform
- System-level virtualization does offer interesting features
  - Application isolation within virtual machines (users can do whatever they want)
  - All standard UNIX tools can be used



## **Challenges for the Usage of VSE in Distributed Systems**

- System management
  - How can we deploy & manage both the HostOS and the VMs?
  - How can users specify the system within a VM?
  - How can sysadmin specify their constraints regarding execution environments?
  - How can we switch between different virtualization solutions?
  - How can we switch between a virtual environment and a normal environment (e.g., disk-less, disk-full)?

### **Important lack of tools**

 Selected approach: the definition of Virtual System Environments



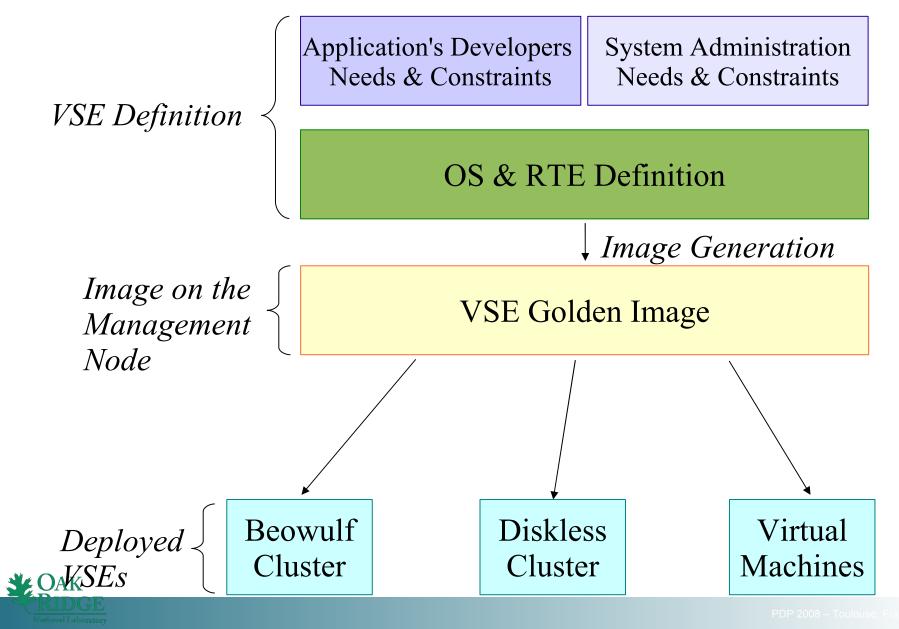
## **OSCAR-V: OSCAR Extension for Virtual** Systems

### Implement VSEs support

- without rewriting everything from scratch
- potential support of all RPM and Debian based Linux distributions
- abstracting existing system-level virtualization solutions
- Provide an integrated solution for:
  - the deployment & management of both HostOSes and VMs
  - the specification of VSE for both the user point-of-view and the sysadmin point-of-view
  - unique interface for the manage of VMs: concept of *profile*
  - possible switch between disk-less, beowulf and partitioned systems (ongoing work)



## **VSE Management - Overview**



## **Abstraction of the System-Level Virtualization Solutions**

- Users do not care about the technical details
  - hide all the configuration details: Virtual Machine Management V2M
  - provide a simple API for the definition of VMs: profiles

### Profile example

```
<?xml version="1.0"?>
<!DOCTYPE profile PUBLIC "" "v3m_profile.dtd">
<profile>
<name>test</name>
<type>Xen</type>
<image size="50">/home/gvallee/temp/v2m/test_xen.img</image>
<nic1>
<type>TUN/TAP</type>
<mac>00:02:03:04:05:06</mac>
</nic1>
```

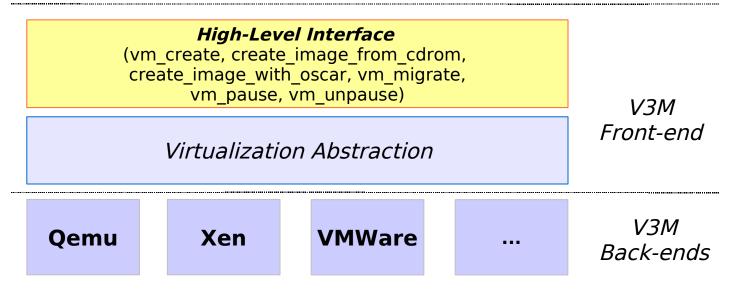


### **V2M - Architecture**

*V2M* (Virtual Machine Management Command Line Interface)

#### *KVMs* (*GUI for Linux - KDE/Qt*)

Applications based on libv3m



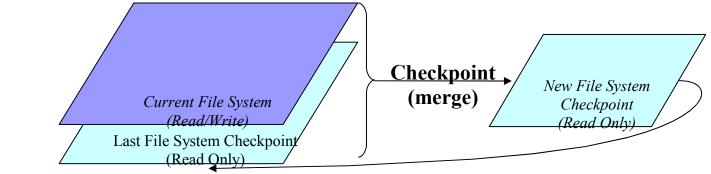


### System-level Virtualization and High Availability



## **VM Checkpoint / Restart**

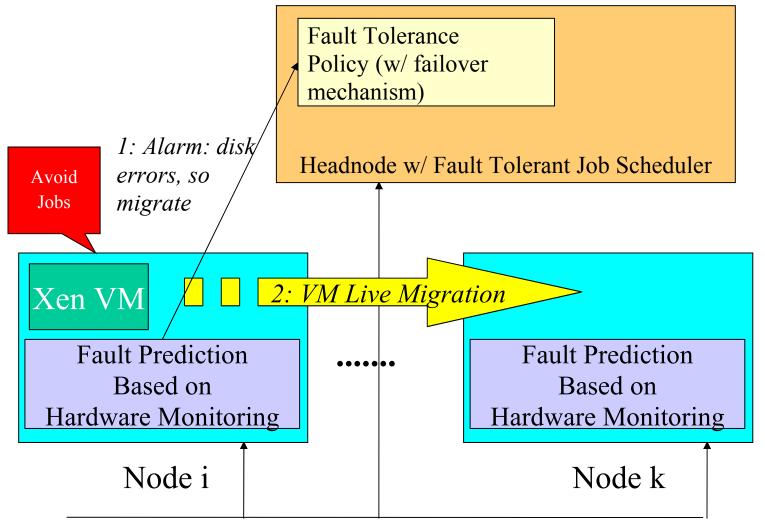
- Already possible to checkpoint the memory (memory dump in a file)
- File system checkpoint/restart



Collaboration with LATech (Box Leangsuksun)



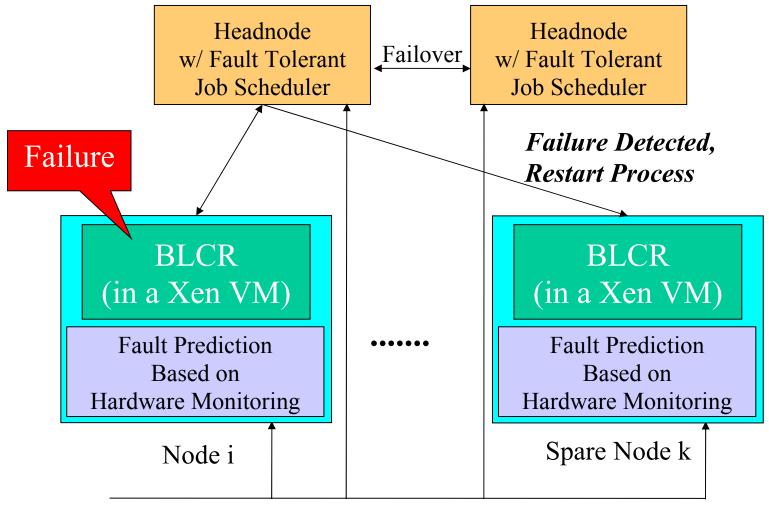
## **Pro-active Fault Tolerance**



Network



## **Reactive Fault Tolerance**



Network



### **System Management & Administration**



### **Resource Management**

### 2 different main challenges

- System partitioning (becomes critical with multi-cores)
- Application deployment

### Virtualization benefits

- easily enable system partitioning
- simplify the resource exposure

### Application/VMs deployment

- on demand
- before application deployment
- Example: Distributed Virtual Clustering (ASU / Cluster Resources Inc.)
  - Extension of MOAB for the virtualization support



### I/O & Storage



## I/O & Storage

 Critical for most of the HPC applications (communications, access to storage on service nodes)

#### Access to the bare hardware

- currently through the Host OS (isolation)
- implies an overhead

#### Possible solutions

- VMM-bypass (direct mapping of resource into VMs)
- Remote Direct Memory Access (RDMA)



## **Impact of Virtualization for HPC**

### Programming paradigm

- Challenges: How to move data to/from the application? How to parallelize applications to hundreds or even thousand of nodes? How to checkpoint/restart applications in order to guarantee resiliency?
- Opportunities: implicit communications move the application to data; change the resource exposure

### Application development

- emulation vs. virtualization; OS adaptation; VSE
- application developers can focus on the science
- System administration
  - separate system administrators and users constraints (VSE)

### Foster research and education

ease research in architecture & operating systems research



## Conclusion

- Virtualization for HPC implies several challenges
  - a hypervisor suitable for HPC (i.e., with a small system footprint)
  - the support of virtual system environments
  - the support of high availability and fault tolerance capabilities
  - the support of advanced resource management capabilities
  - the use of system-level virtualization for resource management
  - the support of efficient I/O mechanisms and storage solutions for virtualized environment.
- No current commercial solution provides such capabilities
- Virtualization solution are still immature for HPC (even if studied since the 70's) and still a lot of research to do



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### **Questions?**

